QOS Management for camera networks

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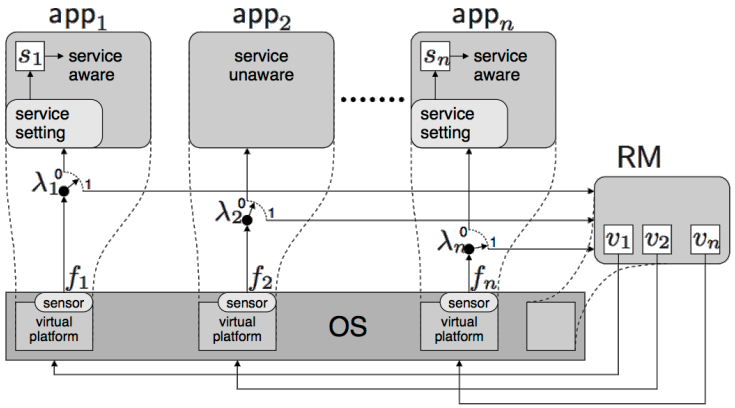
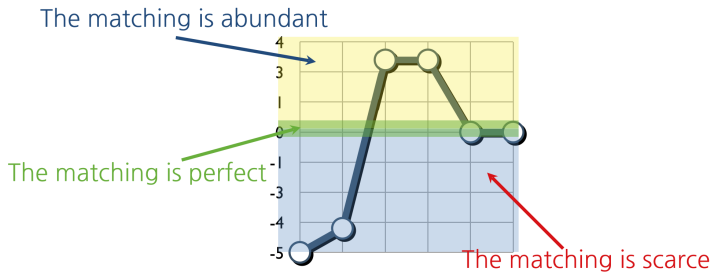
Introduction

- Adaptive resource management for camera networks
- Connected at a local level through Ethernet connection
- Contracts define the resolution and frame rates of the cameras
- Based on previous work by Martina on multicore systems

Resource Manager for Multicore Systems

- Dynamically allocates resources and applications choose their service levels
  - λ_i = 0 means that application adjusts
  - λ_i = 1 means that the manager adjusts
- Decides on the allocated amount depending upon the matching function
  - $f_i < 0$ means matching is scarce
  - $f_i > 0$ means matching is abundant
  - $f_i = 0$ means matching is perfect

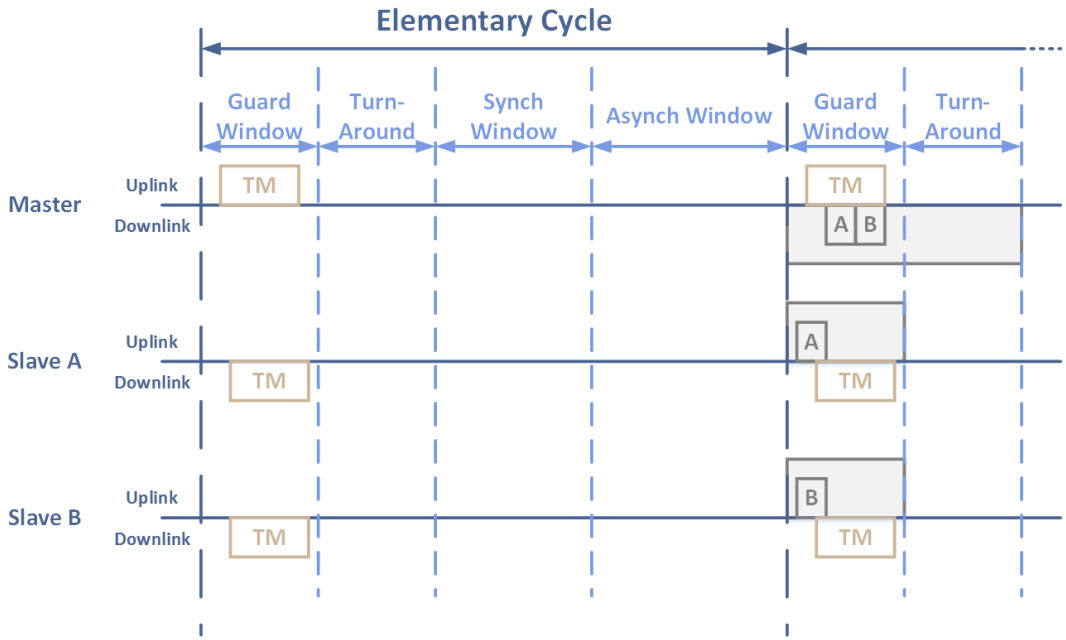




FTT-SE

- Flexible Time Triggered paradigm triggered by a master
- Dynamic reconfiguration of slaves
- Admission control
- Traffic controlled by a master on network
- Switched ethernet, FTT-SE is implemented on COTS switch

Elementary Cycle



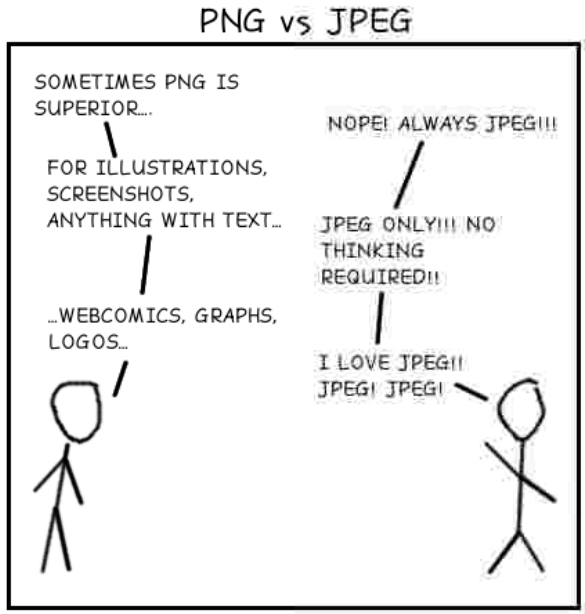
- EC is the traffic schedule in FTT-SE
- The elementary cycle consists of [TM], Async and Sync windows
- Synchronous messages are scheduled by the master.
- Asynchronous messages are sent by the slaves

The Camera Setup

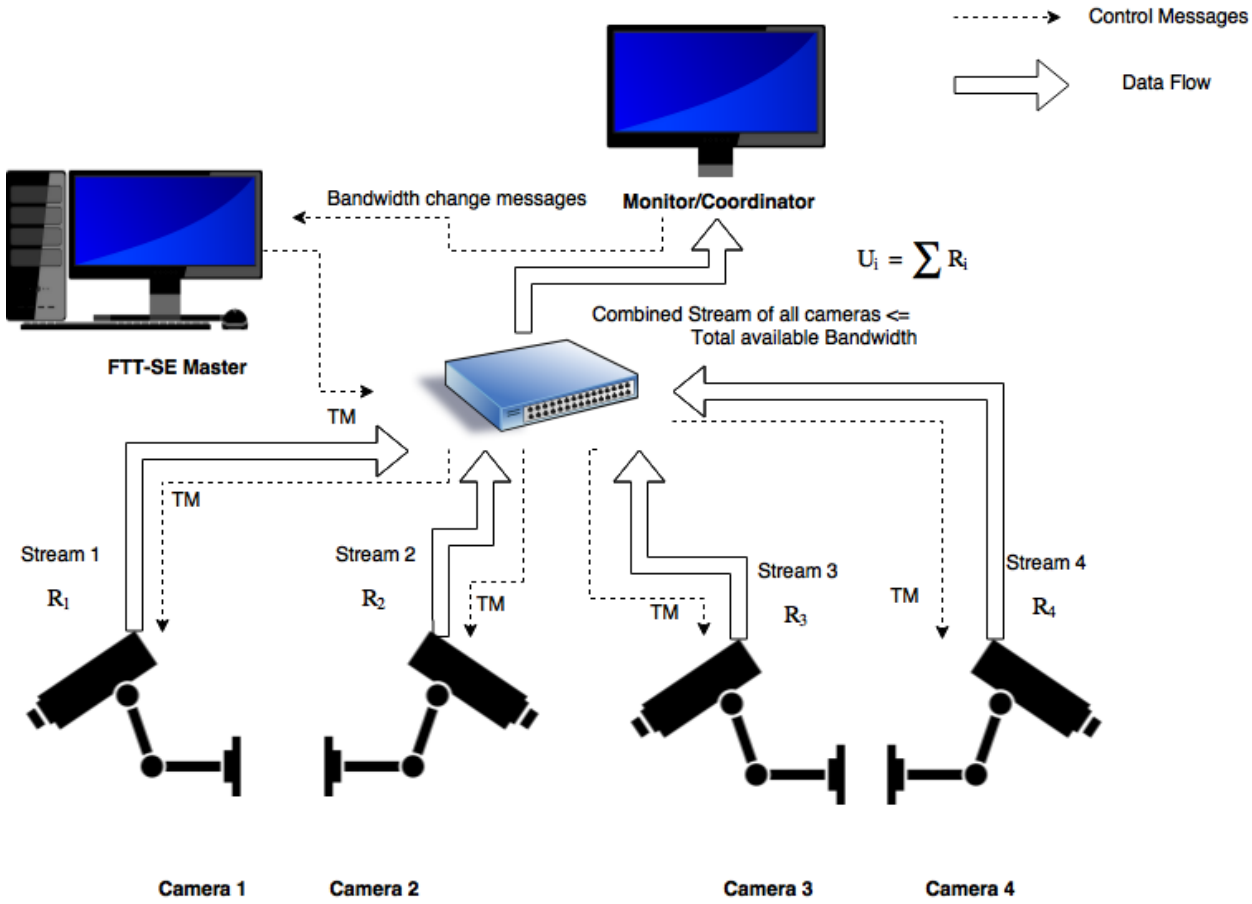
- 4 nodes connected to a switch.
  - One master node, one monitor, and two cameras.
- Master
  - Allocates bandwidths for slaves connected.
  - Notifies BW changes
- Monitor
  - Receives the data/image
  - Monitors BW usage
  - Notifies master for BW changes
- Currently only uses Async messages

A look at the cameras

- Cameras
  - Off the shelf Logitech cameras
  - Captures, encodes and sends the data/image to monitor
  - JPEG compression using OpenCV
  - JPEG is suitable to process and store images
  - Buffer size is varied to fit image



Overview


Control Messages
Data Flow
Monitor/Coordinator
Bandwidth change messages
$U_i = \sum R_i$
Combined Stream of all cameras <=
Total available Bandwidth
FTT-SE Master
TM
TM
Stream 1
R_1
Stream 2
R_2
TM
TM
Stream 3
R_3
TM
Stream 4
R_4
Camera 1
Camera 2
Camera 3
Camera 4

What do they know

- Master
  - Slave channels, Global settings
- Monitor
  - Global settings, Camera contract, Camera state
- Cameras
  - Camera Contract, Camera State
- This allows two control loops. Global QoS and camera
  - Global QoS approach - Game theoretic resource manager
  - On the cameras - PI controller

QoS Approach

- Steady state calculation

$$
\begin{aligned}
B_i^{ss} &= b_i^{ss} B_c \\
b_i^{ss} &= w^p \frac{\rho_i}{\sum_{j=1}^{n} \rho_j} + w^q \frac{q_i}{\sum_{j=1}^{n} q_j^a}
\end{aligned}
$$

- Error function

$$
f_i = \alpha \left(\frac{B_i^{ss} - B_i}{B_i^{ss}} \right)_{[-1\ 1]} + (1-\alpha) \left(\frac{R_i^{ss} - R_i}{B_i^{ss}} \right)_{[-1\ 1]}
$$

- if f_i is less than 0, matching is scarce
- if f_i is greater than 0, matching is abundant
- if f_i is equal to 0, matching is perfect

Game Theoretic Manager

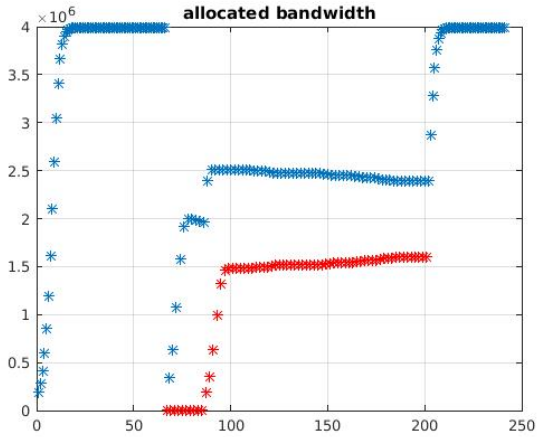
- Present in the Monitor
- Does the following
  1. Measure the current performance $f_i(t)$
  2. Updates the channel as

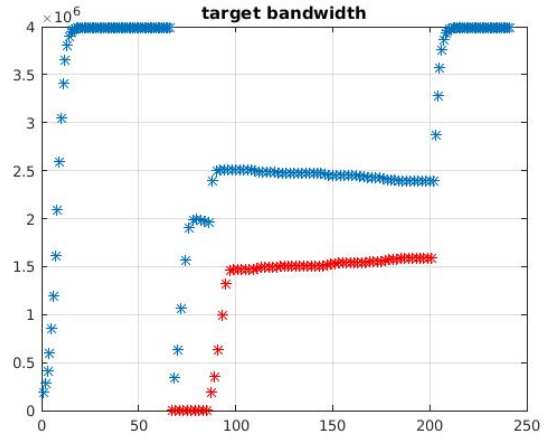
$$bw_i(t+1) = bw_i(t) + \epsilon(-\lambda_i f_i(t) + \sum_{j=1}^{n} \lambda_j f_j(t)\, bw_i(t))$$

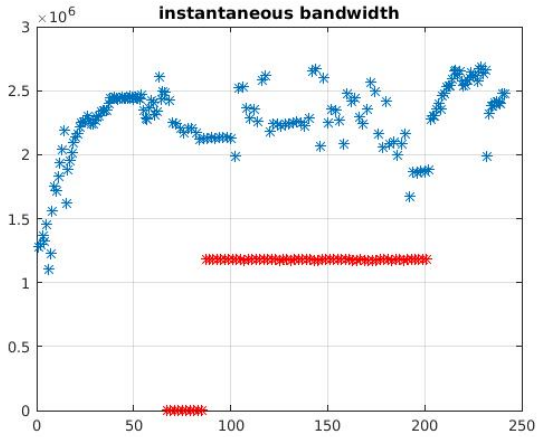
  3. Computes value of bandwidth to allocate

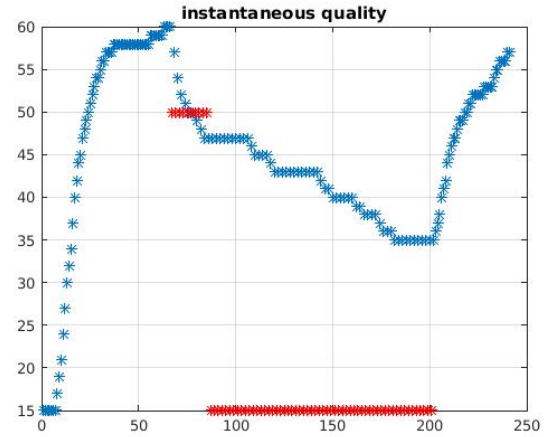
$$B_i(t) = n * bw_i(t+1)$$

Preliminary Results

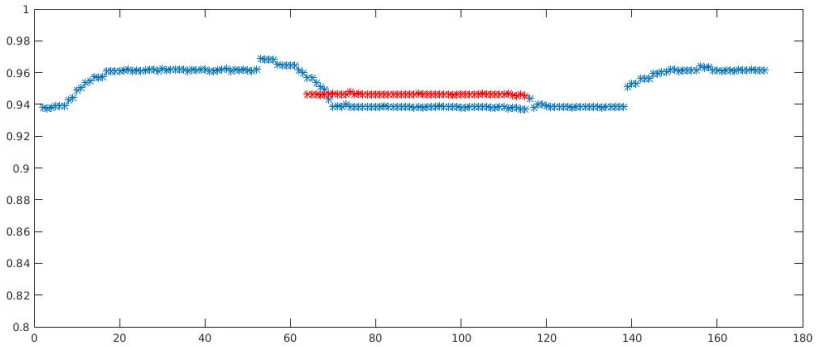


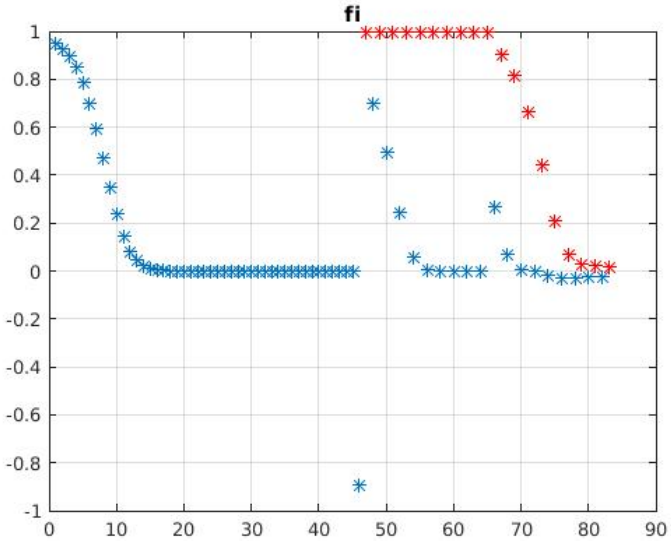






SSIM and f_i


1
0.98
0.96
0.94
0.92
0.9
0.88
0.86
0.84
0.82
0.8
0
20
40
60
80
100
120
140
160
180


fi
1
0.8
0.6
0.4
0.2
0
-0.2
-0.4
-0.6
-0.8
-1
0
10
20
30
40
50
60
70
80
90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AD VT RVMQVE
RVM·CAROLINÆ

Thank you